Taylor Expansion is Game Semantics

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Following the work presented in [BCVA23] on the link between resource λ -calculus and game semantics, we show that the Taylor expansion of a λ -term is isomorphic to its interpretation in pointer concurrent games.

Resource λ -calculus arose from linear logic and differential λ -calculus, and has a strongly finitary behavior. Tsukada and Ong showed in [TO16] that (β -normal, η -long) resource terms correspond to plays of Hyland-Ong games (up to Mellies' homotopy equivalence). These results were extended in [BCVA23], using *pointer concurrent games*, a game model inspired from concurrent games which represents plays quotiented by homotopy. Normal, η -long resource terms are isomorphic to *augmentations* (canonical representatives of Hyland-Ong plays up to homotopy), and that isomorphism is compatible with the β reduction, giving us the following diagram:

$$\begin{array}{ccc} s & \longmapsto & \mathcal{N}(s) \\ \llbracket - \rrbracket & & \downarrow & \downarrow \\ \llbracket s \rrbracket & = & \llbracket \mathcal{N}(s) \rrbracket \end{array}$$

where s is an η -long resource term, \mathcal{N} is the normalisation, and $[\![-]\!]$ is the interpretation in pointer concurrent games.

Taylor expansion translates a λ -term (with possibly infinite behavior) to an infinite sum of resource λ -terms. In this work in progress, we extend the previous isomorphism to show that game semantics is compatible with Taylor expansion in the following sense:

$$\begin{array}{cccc} M & \longmapsto & \mathcal{T}(M) & \longmapsto & \mathcal{N}(\mathcal{T}(M)) \\ \mathbb{I} & & \mathbb{I} & & \mathbb{I} \\ \mathbb{I} & & & \mathbb{I} \\ \mathbb{I} & & & \mathbb{I} \\ \mathbb{I} & & & \mathbb{I} \end{array} \end{array}$$

where M is a λ -term and \mathcal{T} is the Taylor expansion. To do so, we define a Taylor expansion sending simply-typed λ -terms to terms of the simply-typed, η -long resource calculus, and we show that (1) commutes – then (2) is obtained from (*).

References

- [BCVA23] Lison Blondeau-Patissier, Pierre Clairambault, and Lionel Vaux Auclair. Strategies as Resource Terms, and Their Categorical Semantics. In Marco Gaboardi and Femke van Raamsdonk, editors, 8th International Conference on Formal Structures for Computation and Deduction (FSCD 2023), volume 260 of Leibniz International Proceedings in Informatics (LIPIcs), pages 13:1–13:22, Dagstuhl, Germany, 2023. Schloss Dagstuhl – Leibniz-Zentrum für Informatik.
- [TO16] Takeshi Tsukada and C.-H. Luke Ong. Plays as resource terms via non-idempotent intersection types. In Martin Grohe, Eric Koskinen, and Natarajan Shankar, editors, Proceedings of the 31st Annual ACM/IEEE Symposium on Logic in Computer Science, LICS '16, New York, NY, USA, July 5-8, 2016, pages 237–246. ACM, 2016.