## Introduction to Data-Driven Dependency Parsing

Introductory Course, ESSLLI 2007

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## **Overview of the Course**

- Dependency parsing (Joakim)
- Machine learning methods (Ryan)
- Transition-based models (Joakim)
- Graph-based models (Ryan)
- Loose ends (Joakim, Ryan):
  - Other approaches
  - Empirical results
  - Available software

## **Notation Reminder**

- Sentence  $x = w_0, w_1, \ldots, w_n$ , with  $w_0 = root$
- $L = \{l_1, \ldots, l_{|L|}\}$  set of permissible arc labels
- Let G = (V, A) be a dependency graph for sentence x where:
  - $V = \{0, 1, \dots, n\}$  is the vertex set
  - A is the arc set, i.e., (i, j, k) ∈ A represents a dependency from w<sub>i</sub> to w<sub>j</sub> with label l<sub>k</sub> ∈ L
- By the usual definition, G is a tree

#### **Data-Driven Parsing**

- Goal: Learn a good predictor of dependency graphs
- ▶ Input: x
- Output: dependency graph/tree G
- This lecture:
  - Parameterize parsing by transitions
  - Learn to predict transitions given the input and a history
  - Predict new graphs using deterministic parsing algorithm
- Next lecture:
  - Parameterize parsing by dependency arcs
  - Learn to predict entire graphs given the input
  - Predict new graphs using spanning tree algorithms

#### Lecture 3: Outline

- Transition systems
- Deterministic classifier-based models
  - Parsing algorithm
  - Stack-based and list-based transition systems
  - Classifier-based parsing
- Pseudo-projective parsing

#### **Transition Systems**

- ► A transition system for dependency parsing is a quadruple S = (C, T, c<sub>s</sub>, C<sub>t</sub>), where
  - 1. C is a set of configurations, each of which contains a buffer  $\beta$  of (remaining) nodes and a set A of dependency arcs,
  - 2. T is a set of transitions, each of which is a (partial) function  $t: C \rightarrow C$ ,
  - 3.  $c_s$  is an initialization function, mapping a sentence
    - $x = w_0, w_1, \ldots, w_n$  to a configuration with  $\beta = [1, \ldots, n]$ ,
  - 4.  $C_t \subseteq C$  is a set of terminal configurations.
- Note:
  - A configuration represents a parser state.
  - A transition represents a parsing action (parser state update).

### **Transition Sequences**

- Let  $S = (C, T, c_s, C_t)$  be a transition system.
- A transition sequence for a sentence x = w<sub>0</sub>, w<sub>1</sub>,..., w<sub>n</sub> in S is a sequence C<sub>0,m</sub> = (c<sub>0</sub>, c<sub>1</sub>,..., c<sub>m</sub>) of configurations, such that

1. 
$$c_0 = c_s(x)$$

2. 
$$c_m \in C_t$$
,

- 3. for every  $i \ (1 \le i \le m)$ ,  $c_i = t(c_{i-1})$  for some  $t \in T$ .
- ► The parse assigned to x by C<sub>0,m</sub> is the dependency graph G<sub>cm</sub> = ({0, 1, ..., n}, A<sub>cm</sub>), where A<sub>cm</sub> is the set of dependency arcs in c<sub>m</sub>.

#### **Deterministic Parsing**

- An oracle for a transition system S = (C, T, c<sub>s</sub>, C<sub>t</sub>) is a function o : C → T.
- ► Given a transition system S = (C, T, c<sub>s</sub>, C<sub>t</sub>) and an oracle o, deterministic parsing can be achieved by the following simple algorithm:

Parse
$$(x = (w_0, w_1, \dots, w_n))$$
  
1  $c \leftarrow c_s(x)$   
2 while  $c \notin C_t$   
3  $c = [o(c)](c)$   
4 return  $G_c$ 

▶ NB: Oracles can be approximated by classifiers (cf. lecture 2).

#### Stack-Based Transition Systems

- A stack-based configuration for a sentence  $x = w_0, w_1, \ldots, w_n$ is a triple  $c = (\sigma, \beta, A)$ , where
  - 1.  $\sigma$  is a stack of tokens  $i \leq m$  (for some  $m \leq n$ ),
  - 2.  $\beta$  is a buffer of tokens i > m,
  - 3. A is a set of dependency arcs such that  $G = (\{0, 1, \dots, n\}, A)$ is a dependency graph for x.
- A stack-based transition system is a quadruple S

$$C = (C, T, c_s, C_t)$$
, where

1. C is the set of all stack-based configurations,

2. 
$$c_s(x = w_0, w_1, \dots, w_n) = ([0], [1, \dots, n], \emptyset),$$

3. T is a set of transitions, each of which is a function  $t: C \to C$ ,

4. 
$$C_t = \{c \in C | c = (\sigma, [], A)\}.$$

Notation:

•  $\sigma | i =$  stack with top i (| left-associative)

•  $i|\beta$  = buffer with next token *i* (| right-associative)

## Shift-Reduce Dependency Parsing

- ► Transitions:
  - ► Left-Arc<sub>k</sub>:

 $(\sigma|i,j|\beta,A) \Rightarrow (\sigma,j|\beta,A \cup \{(j,i,k)\})$ 

▶ Right-Arc<sub>k</sub>:

$$(\sigma|i,j|\beta,A) \Rightarrow (\sigma,i|\beta,A\cup\{(i,j,k)\})$$

► Shift:

$$(\sigma, i | \beta, A) \Rightarrow (\sigma | i, \beta, A)$$

Preconditions:

 $[root_0]_{\sigma}$  [Economic<sub>1</sub> news<sub>2</sub> had<sub>3</sub> little<sub>4</sub> effect<sub>5</sub> on<sub>6</sub> financial<sub>7</sub> markets<sub>8</sub> .9]<sub> $\beta$ </sub>

 $[root_0 \ Economic_1]_{\sigma} \ [news_2 \ had_3 \ little_4 \ effect_5 \ on_6 \ financial_7 \ markets_8 \ .9]_{\beta}$ 

#### nmod

 $[root_0]_{\sigma}$  Economic<sub>1</sub> [news<sub>2</sub> had<sub>3</sub> little<sub>4</sub> effect<sub>5</sub> on<sub>6</sub> financial<sub>7</sub> markets<sub>8</sub> .<sub>9</sub>]<sub> $\beta$ </sub>

Left-Arcnmod

#### nmod

 $[root_0 \ Economic_1 \ news_2]_{\sigma} \ [had_3 \ little_4 \ effect_5 \ on_6 \ financial_7 \ markets_8 \ .9]_{\beta}$ 

# $[root_0]_{\sigma} \ Economic_1 \ news_2 \ [had_3 \ little_4 \ effect_5 \ on_6 \ financial_7 \ markets_8 \ .9]_{\beta}$

Left-Arc<sub>sbj</sub>

# $[root_0 \ Economic_1 \ news_2 \ had_3]_{\sigma} \ [little_4 \ effect_5 \ on_6 \ financial_7 \ markets_8 \ .9]_{\beta}$

# $[root_0 \ Economic_1 \ news_2 \ had_3 \ little_4]_{\sigma} \ [effect_5 \ on_6 \ financial_7 \ markets_8 \ .9]_{\beta}$



Left-Arcnmod









Left-Arcnmod

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Right-Arc<sub>pc</sub>



Right-Arcnmod

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Right-Arcobj



#### Right-Arcpred



Right-Arc<sub>p</sub>



#### **Theoretical Results**

- Complexity:
  - Deterministic shift-reduce parsing has time and space complexity O(n), where n is the length of the input sentence.
- Correctness:
  - ▶ For every transition sequence  $C_{0,m}$ ,  $G_{c_m}$  is a projective dependency forest (soundness).
  - ▶ For every projective dependency forest *G*, there is a transition sequence  $C_{0,m}$  such that  $G_{c_m} = G$  (completeness).
- Note:
  - A dependency forest is (here) a dependency graph satisfying Root, Single-Head, and Acyclicity (but not Connectedness).
  - A dependency forest G = (V, A) can be transformed into a dependency tree by adding arcs of the form (0, i, k) (for some l<sub>k</sub> ∈ L) for every root i ∈ V (i ≠ 0).

## Variations on Shift-Reduce Parsing

- Empty stack initialization:
  - If we can assume that there is only one node *i* such that (0, *i*, *k*) ∈ *A*, then we can reduce ambiguity by starting with an empty stack (and adding the arc (0, *i*, *k*) after termination).
- Iterative parsing [Yamada and Matsumoto 2003]:
  - Same transition system (with empty stack initialization)<sup>1</sup>
  - Given a terminal configuration:
    - $\blacktriangleright (\sigma, [], A) \Longrightarrow ([], \sigma, A)$
    - ▶ Terminate when A has not been modified in the last iteration.
- Modified transition systems:
  - Arc-eager parsing [Nivre 2003]
  - Non-projective parsing [Attardi 2006]

<sup>1</sup>NB: Left-Arc  $\Rightarrow$  Right, Right-Arc  $\Rightarrow$  Left

# **Arc-Eager Parsing**

- Transitions:
  - Left-Arc<sub>k</sub>:

 $(\sigma|i,j|\beta,A) \Rightarrow (\sigma,j|\beta,A \cup \{(j,i,k)\})$ 

► Right-Arc<sub>k</sub>:

 $(\sigma|i,j|\beta,A) \Rightarrow (\sigma|i|j,\beta,A \cup \{(i,j,k)\})$ 

Reduce:

$$(\sigma|i,\beta,A) \Rightarrow (\sigma,\beta,A)$$

Shift:

$$(\sigma, i|\beta, A) \Rightarrow (\sigma|i, \beta, A)$$

- Preconditions:
  - ► Left-Arc<sub>k</sub>:  $\neg[i = 0]$  $\neg \exists i' \exists k'[(i', i, k') \in A]$
  - Right-Arc<sub>k</sub>:

$$\neg \exists i' \exists k' [(i', j, k') \in A]$$

$$\exists i' \exists k' [(i', i, k') \in A]$$

 $[root_0]_{\sigma}$  [Economic<sub>1</sub> news<sub>2</sub> had<sub>3</sub> little<sub>4</sub> effect<sub>5</sub> on<sub>6</sub> financial<sub>7</sub> markets<sub>8</sub> .9]<sub> $\beta$ </sub>

 $[root_0 \ Economic_1]_{\sigma} \ [news_2 \ had_3 \ little_4 \ effect_5 \ on_6 \ financial_7 \ markets_8 \ .9]_{\beta}$ 

#### nmod

 $[root_0]_{\sigma}$  Economic<sub>1</sub> [news<sub>2</sub> had<sub>3</sub> little<sub>4</sub> effect<sub>5</sub> on<sub>6</sub> financial<sub>7</sub> markets<sub>8</sub> .<sub>9</sub>]<sub> $\beta$ </sub>

Left-Arcnmod

#### nmod

 $[root_0 \ Economic_1 \ news_2]_{\sigma} \ [had_3 \ little_4 \ effect_5 \ on_6 \ financial_7 \ markets_8 \ .9]_{\beta}$ 

# $[root_0]_{\sigma} \ Economic_1 \ news_2 \ [had_3 \ little_4 \ effect_5 \ on_6 \ financial_7 \ markets_8 \ .9]_{\beta}$

Left-Arc<sub>sbj</sub>



Right-Arc<sub>pred</sub>





Left-Arcnmod



Right-Arcobj



Right-Arcnmod

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#### Left-Arcnmod



Right-Arc<sub>pc</sub>











Right-Arc<sub>p</sub>

## **Non-Projective Parsing**

- New transitions:
  - ▶ NP-Left-Arc<sub>k</sub>:

 $(\sigma|i|i',j|\beta,A) \Rightarrow (\sigma|i',j|\beta,A \cup \{(j,i,k)\})$ 

▶ NP-Right-Arc<sub>k</sub>:

$$(\sigma|i|i',j|\beta,A) \Rightarrow (\sigma|i,i'|\beta,A \cup \{(i,j,k)\})$$

Handles most naturally occurring non-projective dependency relations (94% in the Prague Dependency Treebank).



More expressive extensions are possible [Attardi 2006].

## **Comparing Algorithms**

#### Expressivity:

- Arc-standard and arc-eager shift-reduce parsing is limited to projective depedendency graphs.
- Simple extensions can handle a subset of non-projective dependency graphs.
- Complexity:
  - Space complexity is O(n) for all deterministic parsers (even with simple extensions).
  - ► Time complexity is O(n) for single-pass parsers, O(n<sup>2</sup>) for iterative parsers.
- More complex extensions to handle non-projective dependency graphs will affect time complexity.

#### **List-Based Transition Systems**

- A list-based configuration for a sentence x = w<sub>0</sub>, w<sub>1</sub>,..., w<sub>n</sub> is a quadruple c = (λ<sub>1</sub>, λ<sub>2</sub>, β, A), where
  - 1.  $\lambda_1$  is a list of tokens  $i_1 \leq m_1$  (for some  $m_1 \leq n$ ),
  - 2.  $\lambda_2$  is a list of tokens  $i_2 \leq m_2$  (for some  $m_2$ ,  $m_1 < m_2 \leq n$ ),
  - 3.  $\beta$  is a buffer of tokens  $j > m_2$ ,
  - A is a set of dependency arcs such that G = ({0, 1, ..., n}, A) is a dependency graph for x.
- A list-based transition system is a quadruple

$$\mathcal{S} = (\mathcal{C}, \mathcal{T}, \mathit{c_s}, \mathcal{C_t})$$
, where

1. C is the set of all list-based configurations,

2. 
$$c_s(x = w_0, w_1, \dots, w_n) = ([0], [], [1, \dots, n], \emptyset),$$

3. T is a set of transitions, each of which is a function  $t: C \rightarrow C$ ,

4. 
$$C_t = \{ c \in C | c = (\lambda_1, \lambda_2, [], A) \}.$$

Notation:

•  $\lambda_1 | i =$  list with head *i* and tail  $\lambda_1$  (| left-associative)

•  $i|\lambda_2 = i$  and tail  $\lambda_2$  (| right-associative)

#### **Non-Projective Parsing**

- ► Transitions:
  - ► Left-Arc<sub>k</sub>:

 $(\lambda_1|i,\lambda_2,j|\beta,A) \Rightarrow (\lambda_1,i|\lambda_2,j|\beta,A\cup\{(j,i,k)\})$ 

▶ Right-Arc<sub>k</sub>:

 $(\lambda_1|i,\lambda_2,j|\beta,A) \Rightarrow (\lambda_1,i|\lambda_2,j|\beta,A\cup\{(i,j,k)\})$ 

► No-Arc:

$$(\lambda_1|i,\lambda_2,\beta,A) \Rightarrow (\lambda_1,i|\lambda_2,\beta,A)$$

$$(\lambda_1, \lambda_2, i|\beta, A) \Rightarrow (\lambda_1.\lambda_2|i, [], \beta, A)$$

- Preconditions:
  - Left-Arc:

## **Projective Parsing**

- Transitions:
  - ► Left-Arc<sub>k</sub>:

 $(\lambda_1|i,\lambda_2,j|\beta,A) \Rightarrow (\lambda_1,\lambda_2,j|\beta,A\cup\{(j,i,k)\})$ 

▶ Right-Arc<sub>k</sub>:

 $(\lambda_1|i,\lambda_2,j|\beta,A) \Rightarrow (\lambda_1|i|j,[],\beta,A\cup\{(i,k,j)\})$ 

► No-Arc:

 $(\lambda_1|i,\lambda_2,\beta,A) \Rightarrow (\lambda_1,i|\lambda_2,\beta,A)$ 

Shift:

$$(\lambda_1, \lambda_2, i|\beta, A) \Rightarrow (\lambda_1.\lambda_2|i, [], \beta, A)$$

Preconditions:

- ► Left-Arc:  $\neg[i = 0]$  $\neg \exists i' \exists k'[(i', k', i) \in A]$
- Right-Arc:

$$\neg \exists i' \exists k' [(i', k', j) \in A]$$

No-Arc:  $\exists i' \exists k[(i', k, i) \in A]$ 

#### **Theoretical Results**

#### Complexity:

- Deterministic list-based parsing has time complexity O(n<sup>2</sup>) and space complexity O(n), where n is the length of the input sentence.
- Correctness:
  - ▶ For every transition sequence C<sub>0,m</sub>, G<sub>cm</sub> is a (projective) dependency forest (soundness).
  - ▶ For every (projective) dependency forest G, there is a transition sequence  $C_{0,m}$  such that  $G_{c_m} = G$  (completeness).

## **Classifier-Based Parsing**

Data-driven deterministic parsing:

- Deterministic parsing requires an oracle.
- An oracle can be approximated by a classifier.
- A classifier can be trained using treebank data.
- Learning problem:
  - Approximate a function from configurations (represented by feature vectors) to transitions, given a training set of (gold standard) transition sequences.
  - ► Three issues:
    - How do we represent configurations by feature vectors?
    - How do we derive training data from treebanks?
    - How do we learn classifiers?

#### **Feature Representations**

- ► A feature representation f(c) of a configuration c is a vector of simple features f<sub>i</sub>(c).
- Typical features are defined in terms of attributes of nodes in the dependency graph.
  - ► Nodes:
    - Target nodes (top of  $\sigma$ , head of  $\lambda_1$ ,  $\lambda_2$ ,  $\beta$ )
    - Linear context (neighbors in  $\sigma$ ,  $\lambda_1$ ,  $\lambda_2$ , or  $\beta$ )
    - Structural context (parents, children, siblings given A)
  - Attributes:
    - Word form (and/or lemma)
    - Part-of-speech (and morpho-syntactic features)
    - Dependency type (if labeled)
    - Distance (between target tokens)

#### A Typical Model [Nivre et al. 2006]



## **Training Data**

- Training instances have the form  $(\mathbf{f}(c), t)$ , where
  - 1. f(c) is a feature representation of a configuration c,
  - 2. *t* is the correct transition out of *c* (i.e., o(c) = t).
- Given a dependency treebank, we can sample the oracle function o as follows:
  - ▶ For each sentence x with (gold standard) dependency graph  $G_x$ , we construct a transition sequence  $C_{0,m} = (c_0, c_1, ..., c_m)$  such that
    - 1.  $c_0 = c_s(x)$ , 2.  $G_{c_m} = G_x$ ,
  - For each configuration c<sub>i</sub>(i < m), we construct a training instance (f(c<sub>i</sub>), t<sub>i</sub>), where t<sub>i</sub>(c<sub>i</sub>) = c<sub>i+1</sub>.

## Learning Classifiers

- Learning methods:
  - Support vector machines (SVM)
    - [Kudo and Matsumoto 2002, Yamada and Matsumoto 2003, Isozaki et al. 2004, Cheng et al. 2004, Nivre et al. 2006]
      - Polynomial kernel ( $d \ge 2$ )
      - Different techniques for multiclass classification
      - Training efficiency problematic for large data sets
  - Memory-based learning (MBL)
    - [Nivre et al. 2004, Nivre and Scholz 2004, Attardi 2006]
      - k-NN classification
      - Different distance functions
      - Parsing efficiency problematic for large data sets
  - Maximum entropy modeling (MaxEnt) [Cheng et al. 2005, Attardi 2006]
    - Extremely efficient parsing
    - Slightly less accurate

 Technique for non-projective dependency parsing with a data-driven projective parser [Nivre and Nilsson 2005].

Four steps:

- 1. Projectivize dependency graphs in training data, encoding information about transformations in augmented arc labels.
- 2. Train projective parser (as usual).
- 3. Parse new sentences using projective parser (as usual).
- 4. Deprojectivize output dependency graphs by heuristic transformations guided by augmented arc labels.

- Projectivize training data:
  - Projective head nearest permissible ancestor of real head
  - Arc label extended with dependency type of real head



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  - Projective head nearest permissible ancestor of real head
  - Arc label extended with dependency type of real head



- Deprojectivize parser output:
  - Top-down, breadth-first search for real head
  - Search constrained by extended arc label



- Deprojectivize parser output:
  - Top-down, breadth-first search for real head
  - Search constrained by extended arc label



## Summary – Transition-based Methods

- Transition systems
- Deterministic classifier-based parsing
  - Parsing algorithm
  - Stack-based and list-based transitions systems
  - Classifier-based parsing
- Pseudo-projective parsing

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