## On the Longest Flip Sequence to Untangle Segments in the Plane

Guilherme da Fonseca - Aix-Marseille Université and LIS, France<br>Yan Gerard - Université Clermont Auvergne and LIMOS, France Bastien Rivier - Université Clermont Auvergne and LIMOS, France

Section 1

## Introduction

## Motivation: Untangling TSP Tours

- 2d Euclidean TSP (NP-hard):

Input: A set of $n$ points called cities.
Output: The shortest tour (polygon whose vertices are the cities).
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## A Potential Argument

■ An infinite flip sequence?

- Measuring progress with a potential,
i.e., an integer function on tours which is:
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$$
1 \leq \underbrace{\Phi_{\text {rank }} \text { when sorted by length }(T)}_{\text {potential of the tour } T} \leq n!
$$



## The Longest Flip Sequences

- The deletion choice may impact the number of flips.
- We know of no clever way to choose.
- $\mathbf{D}(n)$ : number of flins in the longest flip sequences.
- We want bounds on $\mathbf{D}(n)$.


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## Previous Bounds on the Longest Flip Sequences



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## Section 2

## Convex Proof

$n^{2} \preccurlyeq \mathbf{D}_{\text {convex }}(n) \preccurlyeq n^{2}$
$n^{2} \preccurlyeq \mathbf{D}(n, t) \preccurlyeq n^{2} t$
$n^{2} \preccurlyeq \quad \mathbf{D}(n)$
$\preccurlyeq n^{3}$

- $\Phi_{\text {crossings }}(T)$ : number of crossings in the tour $T$.

■ $\Phi_{\text {crossings }}=O\left(n^{2}\right)$

- $\Phi_{\text {crossings }}$ decreases at each flip:



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## From Segments to Lines

## Introduction

Convex Proof

- $\Phi_{\text {crossings }}$ may not decrease at each flip:
- Idea: consider line-segment crossings instead.
- $L$ : lines through two cities.
- $\Phi_{\ell}(T)$ : number of crossings with a line $\ell$ in the tour $I$
- $\Phi_{L}=\sum_{\ell \in L} \Phi$
- $O\left(n^{2}\right)$ lines, $O(n)$ crossings per line $\Longrightarrow \Phi_{L}=O\left(n^{3}\right)$
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## What Is a Crossing

- A single point intersection between a line and a segment is a crossing if it is not an endpoint of the segment.


## Section 4

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n^{2} & \preccurlyeq & \mathbf{D}(n) &
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## Near Convex Sets

■ Near Convex sets: the $n$ points are convex except $t$ of them.

$$
n=9 \quad t=3
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1980 Proof
New Proof Near Convex Mixed Potential Bounded Decreasing Conclusion


- Case 1. If $\Phi_{\text {crossings }}$ decreases, then so does $\Phi$ (because $\Phi_{L^{\prime}}$ does not increase) $\checkmark$
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## Does $\Phi$ Decreases? Yes!



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## Section 5

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## From Tours to Segments

- Being a tour is not used in the proofs.


■ $\mathbf{D}, \mathbf{D}_{\mathrm{RB}}, \mathbf{D}_{\mathrm{TSP}}$ are the same asymptotically.

## Conclusion



Thank you!

## Reductions

- $2 \mathbf{D}(n) \leq \mathbf{D}_{\mathrm{RB}}(2 n) \leq \mathbf{D}(2 n)$
- $2 \mathbf{D}_{\mathrm{RB}}(n) \leq \mathbf{D}_{\mathrm{TSP}}(3 n) \leq \mathbf{D}(3 n)$



## Distinct Flips

■ The same pair of segments can be flipped multiple times in the same sequence.

- Counting distinct flips means that we do not count this multiplicity.
- A balancing argument:
- There are $O\left(\frac{n^{3}}{k}\right)$ flips decreasing $\Phi_{L}$ by at least $k$.
- There are $O\left(n^{2} k^{2}\right)$ flips decreasing $\Phi_{L}$ by less than $k$ :

■ We enumerate them by sweeping a line. o

- We choose $k=n^{1 / 3}$.


